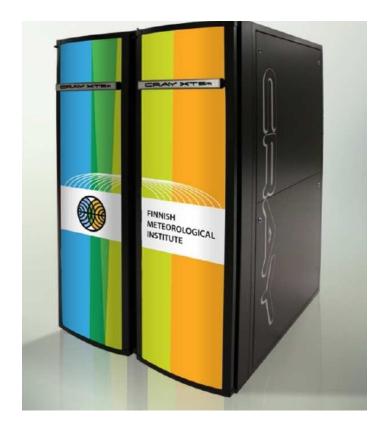
AROME Forecast migration & optimization at FMI



April 13-16, 2010 All Staff Meeting (ASM) Krakow, Poland

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Outline

Introduction

- Fundamental performance problems
- Migration status & some results
- Issues requiring urgent attention
- Good news for benchmarking
- Conclusions





Introduction

FMI has acquired two Cray XT5m clusters with 1968 cores in each

- AROME Forecast migration
 - Harmonie version **35h1.2** (tag 6935) copied to the local FMI-branch
 - Over 3.3 million lines of code (Fortran+C), over 9000 source objects
- CSC IT Center for Science Ltd., (Espoo, FI) was responsible for migrating & optimizing AROME to FMI's Cray XT5



Introduction (cont'd)

AROME model resolution over Finland : 300 x 600 x L40 (2,5km)

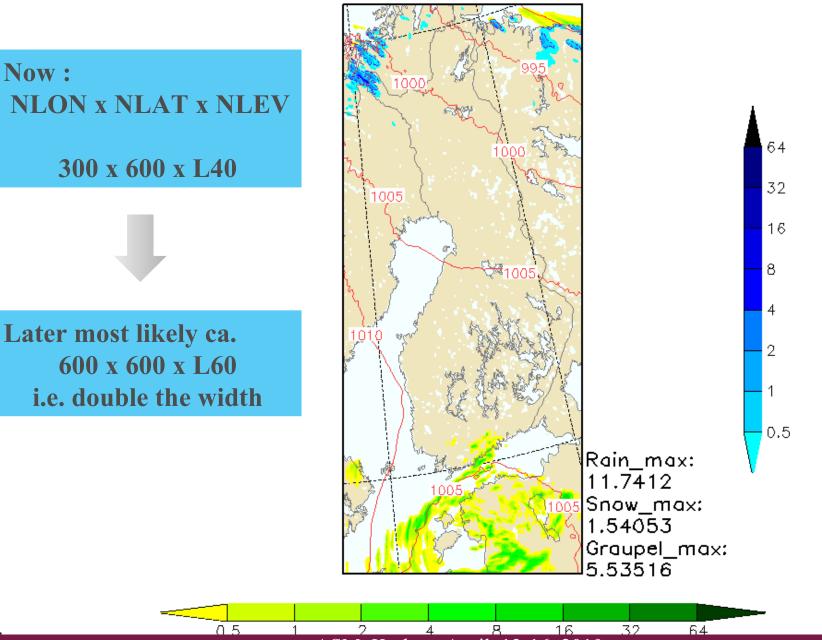
- 24h forecast with radiation calculation every 15th step (every 15 min)
- Full output every 15th step (both Atmos & SURFEX data)
- Boundary input every 60th step

> 24h AROME FC *previously* took 2-3hours on SGI Altix ~ 100 cores

- Was run twice a day in a non time-critical manner
- The near future targets for FMI's AROME on Cray XT5m are
 - Repeat 24h Forecast every 3hrs in a production mode
 - Increase computational area & levels to ca. 600 x 600 x L60 (2,5km)



AROME 090CT2009 00 UTC Forecast. Precipitation [mm 1h⁻¹] 090CT2009 07:00 UTC (aro33h1,2.5km)



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About FMI's Cray XT5m

- Introduction to the system
 - 2 identical clusters, 17.3 TFlop/s peak for each, ca. 35TF total
 - Hex-core AMD Opteron 2.2GHz Istanbul chip

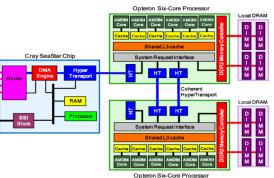
12 (= 2 x 6) cores in a shared memory node
8.8 GFlop/s peak per core, 105.6 Gflop/s peak per node
164 nodes x 12 cores = 1968 cores per each cluster
16 GB shared memory per node (~1.3GB per core)

• 2D-torus SeaStar-1 interconnection network



About FMI's Cray XT5m (cont'd)

- Local Lustre file-system on each cluster
 - 2 X 60TB raw == 2 X 43TB formatted
- 1st cluster was installed in Sep'09, and 2nd late Oct'2009
- Performance acceptance completed on 4th of Nov'2009

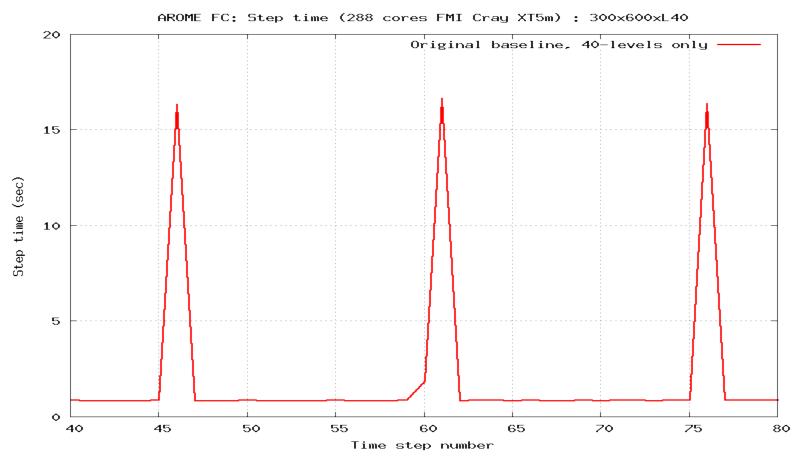


2D Torus Interconned Vetwork Por

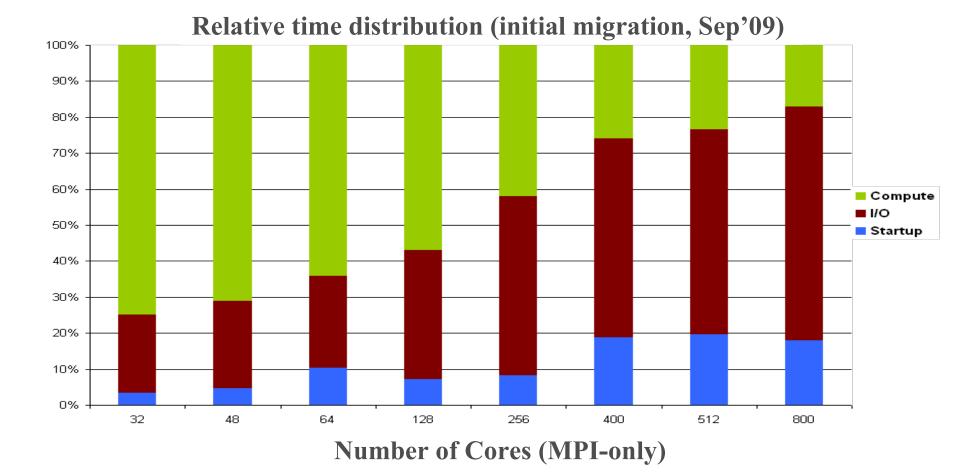




The fundamental performance problem in AROME







Courtesy & thanks to: Markku Kangas, FMI

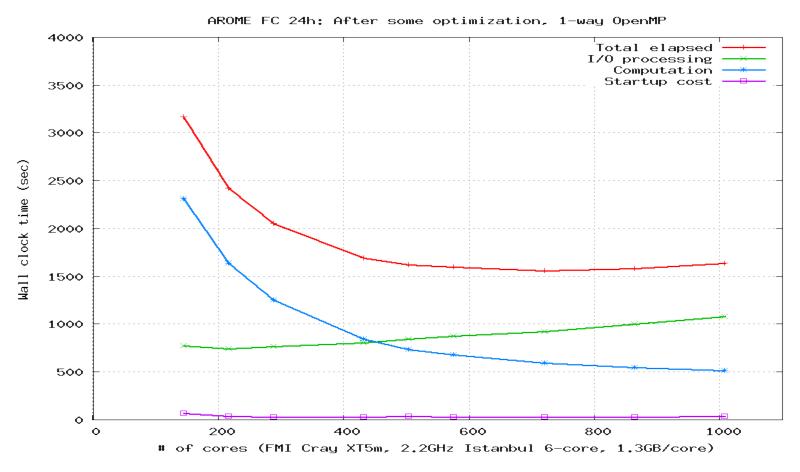


Total **baseline** time = Startup + Comp + I/O

AROME FC 24h: Baseline, no OpenMP Total elapsed I/O processing Computation Startup cost Wall clock time (sec) # of cores (FMI Cray XT5m, 2.2GHz Istanbul 6-core, 1.3GB/core)

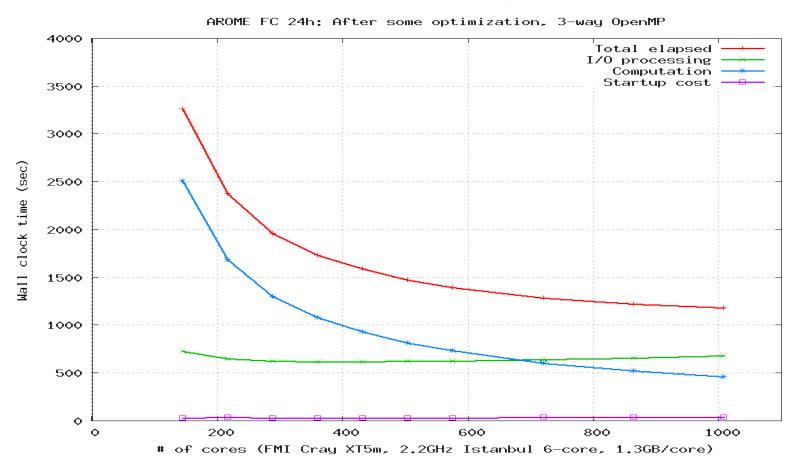


After some optimization : 1-way OpenMP mode





After some optimization : 3-way OpenMP mode





Status of AROME migration to Cray XT5m

AROME Forecast code

- Was difficult to run with > 64 cores due to msg flooding Using synchronized MPI_Ssend and later MPI_gathery in diwrgrid*.F90
- Now works well with MPI + OpenMP > 1700 cores
- Output interval 15min in 24h forecast : 53GB / 24h FC
 I/O processing costs are huge (data gathering, not the I/O itself as such)

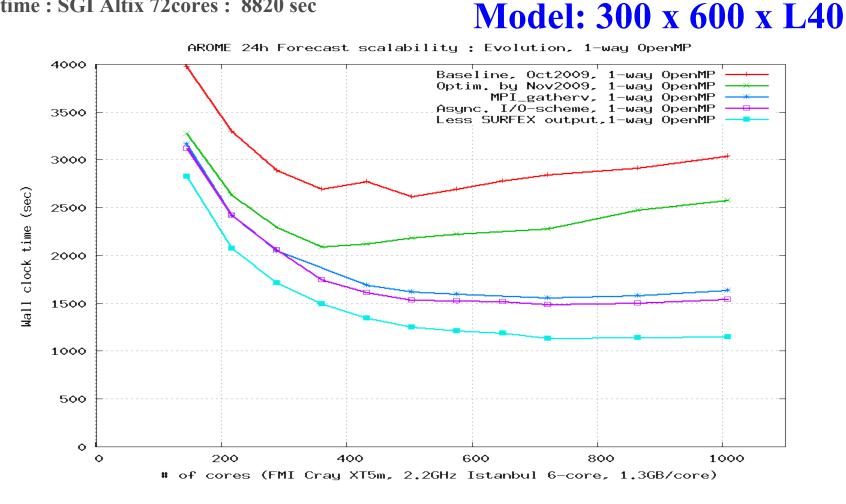


Status of AROME migration ...

- > AROME Forecast code (cont'd)
 - Optimization by looking at the DrHook profiles Savings of 600MB of memory per MPI-task : message passing buffers Weak async I/O scheme (SAMIO) : Comp + I/O cores separated Increase output interval of SURFEX-data from 15min to 1h
- This part was about 2 person months of work
 - Now runs comfortably using 1000 cores with around 10X faster than the reference : runtime under 15-20 minutes



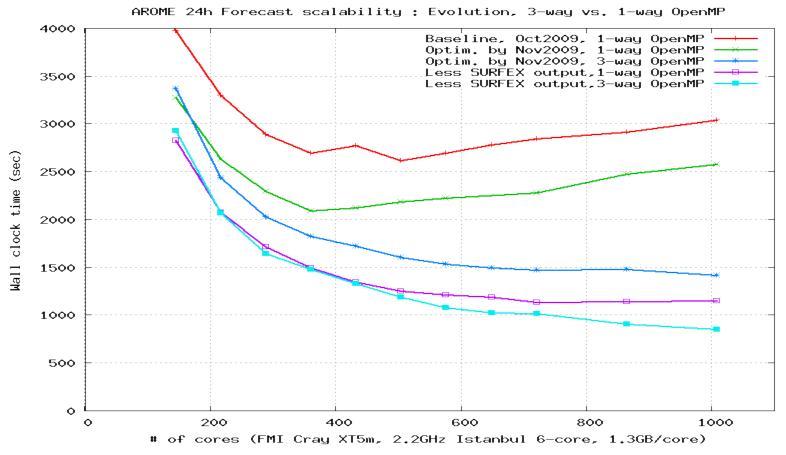
Ref.time : SGI Altix 72cores : 8820 sec





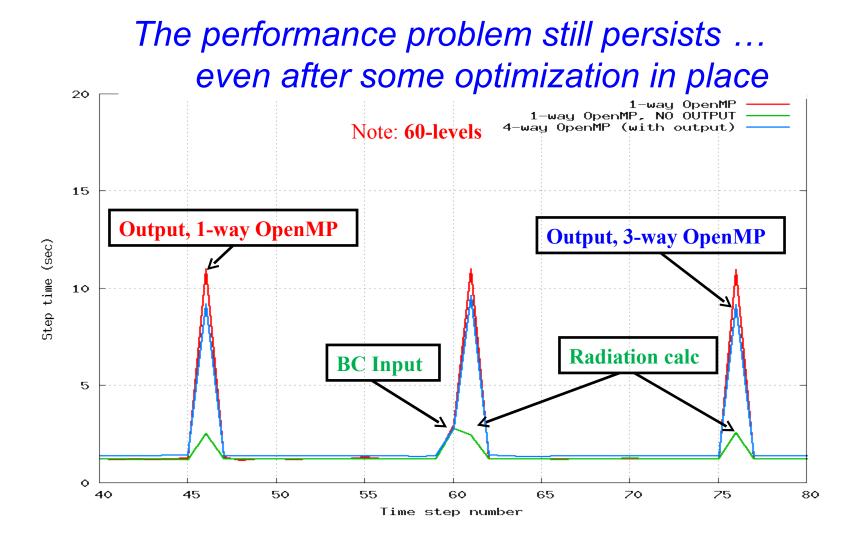
Ref.time : SGI Altix 72cores : 8820 sec

Model: 300 x 600 x L40





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Issues requiring urgent attention

SURFEX & Atmospheric field I/O

- Coded very inefficiently, with parallel computers NOT in mind
- Seriously limits scalability of the whole AROME run
- The I/O logic (data distribution + I/O) needs to be rewritten
- Gains ought to be comparable to Hirlam/HGS (40X speedup in I/O)



Issues requiring ...

Startup cost can be excessive

- 200-300 secs for larger models
- The same BIG files read by all MPI-tasks
- NAMELIST-file processing is pretty slow : Should "HIRLAM technique" be used ?





Issues requiring ...

OpenMP shared memory parallelization part

- Current APL_AROME coding is a suboptimal use of OpenMP
- Prevents exploitation of multicore-technologies





Issues requiring ...

The last, but not least :

- AROME hardly works with other compilers than
 Pathscale, IBM XL and Intel ifort v11 (and of course on NEC)
- GNU Fortran, Portland (PGI), Cray's new FTN still problematic
- Are there genuine problems in the AROME code ?
 There seems to be lots of potentially <u>uninitialized</u> variables !!
 Cray X1E at AEMET finds 285 uninitialized variables !!!
- Can be a *show-stopper* in supercomputer ITT & benchmarking



Good news for benchmarking

New gmake-based robust build mechanism : MAKEUP

- Standard gmake-based & simple to use build system for AROME
- Also builds : Auxlibs (GRIBEX, EMOS), GL-tools, Monitor & Oulan
- Handles ODB builds correctly & in a streamlined manner
- Generates dummy-objects automagically
- Tested so far at FMI, AEMET, SMHI, MetNo, DMI & ECMWF
- One of the January/2010 Harmonie working week topics at FMI



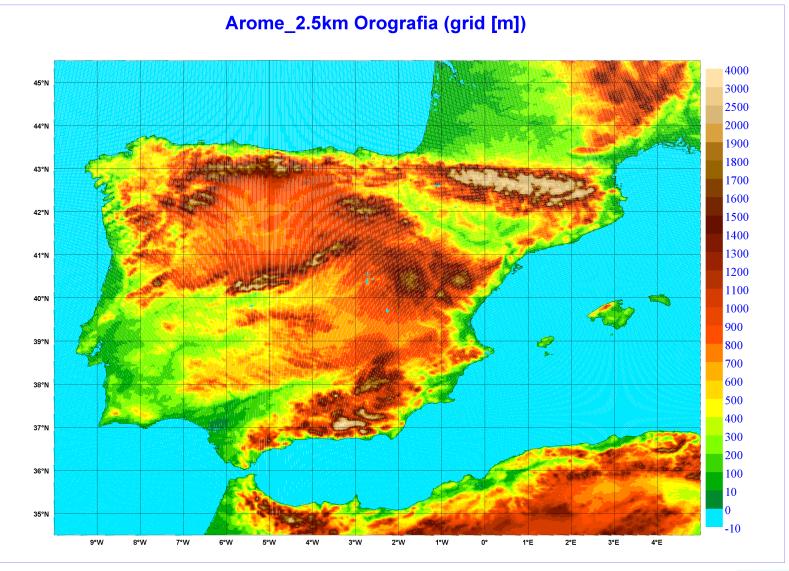
Good news for benchmarking ...

Traceback & performance profile extraction with DrHook

- Useful in debugging & finding locations for crashes
- Produces wall-clock performance profiles cheaply, per MPI-task
- Also supports OpenMP-threads
- Weak asyncronous I/O by using SAMIO
 - SAMIO stands for Storage Access with Multiprocessing I/O
 - A proper solution should be sought after in a collaboration with MF



Benchmark with Iberian peninsula (AEMET)

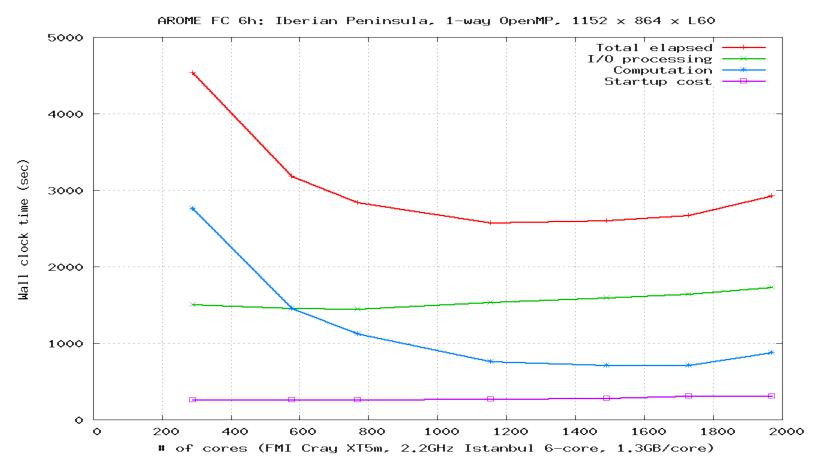




Courtesy & thanks to: Carlos Santos & Jose Antonio Garcia-Moya, AEMET

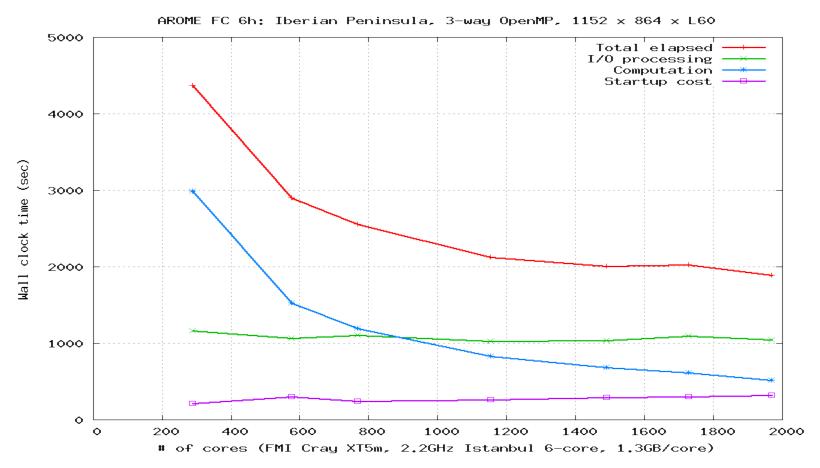
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Iberian peninsula : 1152 x 864 x L60, 1-way OpenMP, 6h FC

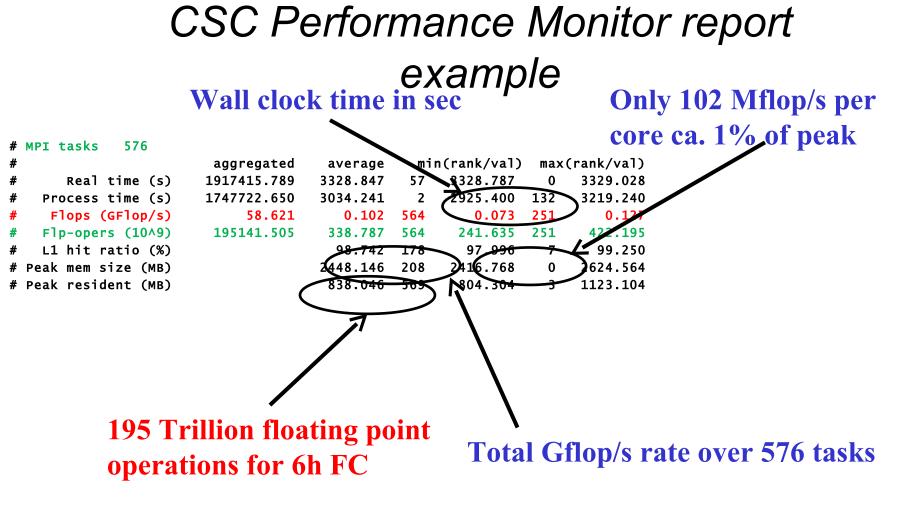




Iberian peninsula : 1152 x 864 x L60, 3-way OpenMP, 6h FC





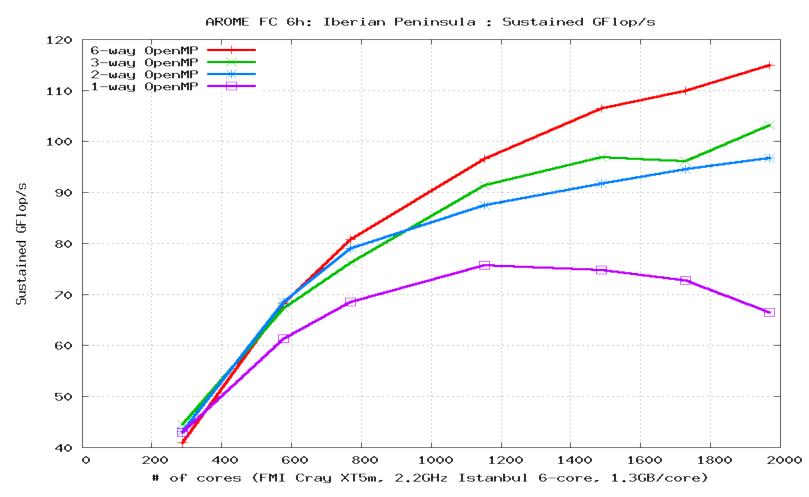


Iberian peninsula 1152 x 864 x L60, 6h forecast 576 MPI-tasks , 1-way OpenMP





Iberian peninsula : Sustained Gflop/s, 6h FC





Conclusions

- AROME FC migration at FMI has demostrated that using hundreds of cores in a mixed MPI +OpenMP parallel environment is a feasible approach
- New MAKEUP build process helps significantly in
 - Program development
 - Standard benchmarking



CSC wonders if AROME could be run further 5X faster ??!

- If I/O problems were fixed
- The causes for high startup cost were fixed
- If OpenMP implementation in APL_AROME was fixed

THANK YOU FOR YOUR ATTENTION !





Extra slides ...

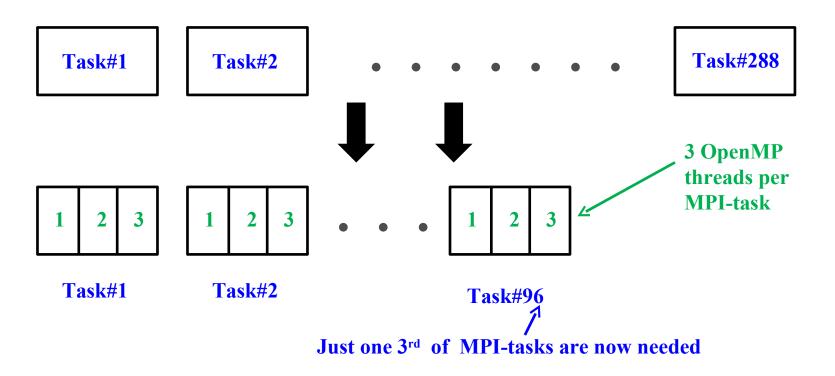




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Using *n-way* OpenMP for AROME FC

- On FMI's Cray XT5m we can use n-way OpenMP for each MPI-task (where n = 1,2,3,4,6 or 12 to keep full cores / node occupancy)
- Use of OpenMP is preferred and gives more bandwidth for AROME
- > For example doing from 1-way to 3-way OpenMP .



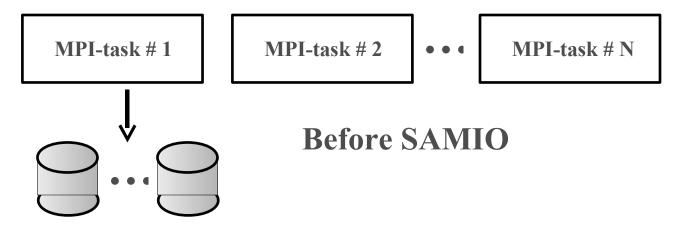


SAMIO

- Storage Access with Multiprocessing I/O
- A "recipe" to replace Fortran I/O with asynchronous I/O with minimal code changes
- Actual I/O performed by separate MPI-tasks, which are not participating in computation at all
- Goals, principles, details …
 - Let computation continue without waiting I/O to complete
 - I/O continues in the "background" by I/O-tasks
 - Upon next output requests currently free I/O-tasks take over
 - An I/O-task responsible for I/O on a given file
 - No MPI I/O, nor split files just low level Unix I/O



Schematic view of SAMIO in AROME





Schematic view of SAMIO in AROME

- Initialize MPI with SAMIO_INIT
- Replace OPEN/READ/WRITE/CLOSE with SAMIO_OPEN/READ/WRITE/CLOSE using (near) original argument lists
- Finish with SAMIO_END
 I/O # 1
 MPI-task# 1
 MPI-task# 2
 MPI-task# N
 I/O # 2
 With SAMIO
 I/O # M



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Some code fractions : OPEN-file

```
SUBROUTINE OPEN_FILE(KNUMER, CDNOMF, CDSTAT, KRECL, ..)
       USE SAMIO_MOD
      IMPLICIT NONE
      LOGICAL LLSAMIO
      LLSAMIO = (SAMIO_MYPE == 1)
      IF (LLSAMIO) THEN
#ifdef USE_SAMIO
           CALL SAMIO_OPEN (UNIT=KNUMER, FILE=CDNOMF, STATUS=CDSTAT,
              ERR=902, FORM='UNFORMATTED', ACCESS='DIRECT', RECL=KRECL,
     &
     &
              IOSTAT=IREP)
           IF (IREP /= 0) GOTO 902
#endif
      ELSE
           OPEN (UNIT=KNUMER, FILE=CDNOMF, STATUS=CDSTAT,
     &
              ERR=902, FORM='UNFORMATTED', ACCESS='DIRECT', RECL=KRECL,
     &
              IOSTAT=IREP)
      ENDIF
```

END SUBROUTINE OPEN_FILE



Some code fractions : WRITE data

```
SUBROUTINE WRITE_CHAR_DATA(KNUMER, KREC, KREOP, CDTAB, ...)
      USE SAMIO_MOD
      IMPLICIT NONE
      CHARACTER CDTAB (JPNXNA*KFACTM)*(JPNCPN)
      LOGICAL LLSAMIO
      LLSAMIO = SAMIO_HAS_OPENED(KNUMER)
      IF (LLSAMIO) THEN
#ifdef USE SAMIO
         CALL SAMIO_WRITE
(UNIT=KNUMER, REC=KREC, ERR=901, IOSTAT=KREP,
     &
       ARRAY=CDTAB)
         IF (KREP /= 0) GOTO 901
#endif
      FI SF
         WRITE (UNIT=KNUMER, REC=KREC, ERR=901, IOSTAT=KREP) CDTAB
      FNDTF
```

```
END SUBROUTINE WRITE_CHAR_DATA
```



. .

What is going on under the hood of SAMIO ?

- The next free I/O-task is chosen during the SAMIO_OPEN and file credentials are sent to that task
- Each SAMIO_WRITE buffers data locally until some user defined buffer size (e.g. 25MB) has been reached
- When cached buffers are full, they are sent to I/O-task using nonblocking (asynchronous) MPI_isend
- SAMIO_READ are processed immediately by requesting read and issuing MPI_recv from I/O-task ...
- ... unless direct access file, in which case this data is available locally – in AROME code at least
- Only SAMIO_CLOSE enforces disk I/O before which an optional (say) big-endian conversion is performed first



Why SAMIO is not the final solution ?

- There are still two memory copies that cost
 - **SAMIO_WRITE** buffering to local cache
 - Local cache is **MPI_Isend** to the remote I/O-task
- Compared to Unix/Linux write caching, there is not much progress – in fact !
- What ought to be done is to send data directly to the I/Oprocesses without intermediary (the 1st computational task currently) and let I/O-tasks to decide how and when and by whom to do the I/O
- Thus there is a need to parallelize properly the data collection steps in AROME I/O, not only the I/O itself

